

# On the global offensive alliance number of a graph

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## Abstract

An offensive alliance in a graph  $\Gamma = (V, E)$  is a set of vertices  $S \subset V$  where for every vertex  $v$  in its boundary it holds that the majority of vertices in  $v$ 's closed neighborhood are in  $S$ . In the case of strong offensive alliance, strict majority is required. An alliance  $S$  is called global if it affects every vertex in  $V \setminus S$ , that is,  $S$  is a dominating set of  $\Gamma$ . The offensive alliance number  $a_o(\Gamma)$  (respectively, strong offensive alliance number  $a_{\delta}(\Gamma)$ ) is the minimum cardinality of an offensive (respectively, strong offensive) alliance in  $\Gamma$ . The global offensive alliance number  $\gamma_o(\Gamma)$  and the global strong offensive alliance number  $\gamma_{\delta}(\Gamma)$  are defined similarly. Clearly,  $a_o(\Gamma) \leq \gamma_o(\Gamma)$  and  $a_{\delta}(\Gamma) \leq \gamma_{\delta}(\Gamma)$ . It was shown in [*Discuss. Math. Graph Theory* **24** (2004), no. 2, 263-275] that  $a_o(\Gamma) \leq \frac{2n}{3}$  and  $a_{\delta}(\Gamma) \leq \frac{5n}{6}$ , where  $n$  denotes the order of  $\Gamma$ . In this paper we obtain several tight bounds on  $\gamma_o(\Gamma)$  and  $\gamma_{\delta}(\Gamma)$  in terms of several parameters of  $\Gamma$ . For instance, we show that  $\frac{2m+n}{3\Delta+1} \leq \gamma_o(\Gamma) \leq \frac{2n}{3}$  and  $\frac{2(m+n)}{3\Delta+2} \leq \gamma_{\delta}(\Gamma) \leq \frac{5n}{6}$ , where  $m$  denotes the

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size of  $\Gamma$  and  $\Delta$  its maximum degree (the last upper bound holds true for all  $\Gamma$  with minimum degree greatest or equal to two).

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## 1 Introduction

The study of defensive alliances in graphs, together with a variety of other kinds of alliances, was introduced in [6]. In the referred paper was initiated the study of the mathematical properties of alliances. In particular, several bounds on the defensive alliance number were given. The particular case of global (strong) defensive alliance was investigated in [4].

The study of offensive alliances was initiated by Favaron et al. in [2] where were derived several bounds on the offensive alliance number and the strong offensive alliance number. On the other hand, in [7] were obtained several tight bounds on different types of alliance numbers of a graph: (global) defensive alliance number, global offensive alliance number and global dual alliance number. In particular, was investigated the relationship between the alliance numbers of a graph and its algebraic connectivity, its spectral radius, and its Laplacian spectral radius. A particular study of the alliance numbers, for the case of planar graphs, can be found in [9]. Moreover, for the study of defensive alliances in the line graph of a simple graph we cite [10].

The aim of this paper is to study mathematical properties of the global offensive alliance number and the global strong offensive alliance number of a graph. We begin by stating some notation and terminology. In this paper  $\Gamma = (V, E)$  denotes a simple graph of order  $n$  and size  $m$ . The degree of a vertex  $v \in V$  will be denoted by  $\delta(v)$ , the minimum degree will be denoted by  $\delta$ , and the maximum degree by  $\Delta$ . The subgraph induced by a set  $S \subset V$  will be denoted by  $\langle S \rangle$ . For a non-empty subset  $S \subset V$ , and a vertex  $v \in V$ , we denote by  $N_S(v)$  the set of neighbors  $v$  has in  $S$ :  $N_S(v) := \{u \in S : u \sim v\}$ . Similarly, we denote by  $N_{V \setminus S}(v)$  the set of neighbors  $v$  has in  $V \setminus S$ :  $N_{V \setminus S}(v) := \{u \in V \setminus S : u \sim v\}$ . The boundary of a set  $S \subset V$  is defined as  $\partial(S) := \bigcup_{v \in S} N_{V \setminus S}(v)$ .

A non-empty set of vertices  $S \subset V$  is called *offensive alliance* if and only if for every  $v \in \partial(S)$ ,  $|N_S(v)| \geq |N_{V \setminus S}(v)| + 1$ . That is, a non-empty set of vertices  $S \subset V$  is called offensive alliance if and only if for every  $v \in \partial(S)$ ,  $2|N_S(v)| \geq \delta(v) + 1$ .

An offensive alliance  $S$  is called *strong* if for every vertex  $v \in \partial(S)$ ,  $|N_S(v)| \geq |N_{V \setminus S}(v)| + 2$ . In other words, an offensive alliance  $S$  is called strong if for every vertex  $v \in \partial(S)$ ,  $2|N_S(v)| \geq \delta(v) + 2$ .

The *offensive alliance number* (respectively, *strong offensive alliance number*), denoted  $a_o(\Gamma)$  (respectively,  $a_\delta(\Gamma)$ ), is defined as the minimum cardinality of an offensive alliance (respectively, strong offensive alliance) in  $\Gamma$ .

A non-empty set of vertices  $S \subset V$  is a *global offensive alliance* if for every vertex  $v \in V \setminus S$ ,  $|N_S(v)| \geq |N_{V \setminus S}(v)| + 1$ . Thus, global offensive alliances are also dominating sets, and one can define the *global offensive alliance number*, denoted  $\gamma_o(\Gamma)$ , to equal the minimum cardinality of a global offensive alliance in  $\Gamma$ . Analogously,  $S \subset V$  is a *global strong offensive alliance* if for every vertex  $v \in V \setminus S$ ,  $|N_S(v)| \geq |N_{V \setminus S}(v)| + 2$ , and the *global strong offensive alliance number*, denoted  $\gamma_\delta(\Gamma)$ , is defined as the minimum cardinality of a global strong offensive alliance in  $\Gamma$ .

In this paper we obtain several tight bounds on  $\gamma_o(\Gamma)$  and  $\gamma_\delta(\Gamma)$  in terms of several parameters of  $\Gamma$ . For instance, we show that

$$\left\lceil \frac{2m+n}{3\Delta+1} \right\rceil \leq \gamma_o(\Gamma) \leq \left\lfloor \frac{2n}{3} \right\rfloor \quad (1)$$

and

$$\left\lceil \frac{2(m+n)}{3\Delta+2} \right\rceil \leq \gamma_\delta(\Gamma) \leq \left\lfloor \frac{5n}{6} \right\rfloor \quad (2)$$

(the last upper bound holds true for all  $\Gamma$  with minimum degree greatest or equal to two).

## 2 Bounding above the global offensive alliance number

It was shown in [2] that the offensive alliance number of a graph of order  $n \geq 2$  is bounded by

$$a_o(\Gamma) \leq \left\lfloor \frac{2n}{3} \right\rfloor, \quad a_o(\Gamma) \leq \left\lfloor \frac{\gamma(\Gamma) + n}{2} \right\rfloor, \quad (3)$$

where  $\gamma(\Gamma)$  denotes the domination number of  $\Gamma$ , and the strong offensive alliance number of a graph of order  $n \geq 3$  is bounded by

$$a_{\delta}(\Gamma) \leq \left\lfloor \frac{5n}{6} \right\rfloor. \quad (4)$$

Clearly,  $a_o(\Gamma) \leq \gamma_o(\Gamma)$  and  $a_{\delta}(\Gamma) \leq \gamma_{\delta}(\Gamma)$ . Now we are going to obtain the above bounds for the case of global alliances.

**Theorem 1.** *For all connected graph  $\Gamma$  of order  $n \geq 2$ ,*

i)  $\gamma_o(\Gamma) \leq \min \left\{ n - \alpha(\Gamma), \left\lfloor \frac{n + \alpha(\Gamma)}{2} \right\rfloor \right\}$ , where  $\alpha(\Gamma)$  denotes the independence number of  $\Gamma$ ;

ii)  $\gamma_o(\Gamma) \leq \left\lfloor \frac{2n}{3} \right\rfloor$ ;

iii)  $\gamma_o(\Gamma) \leq \left\lfloor \frac{\gamma(\Gamma) + n}{2} \right\rfloor$ , where  $\gamma(\Gamma)$  denotes the domination number of  $\Gamma$ ;

iv)  $\gamma_o(\Gamma) \leq \left\lfloor \frac{n(2\mu - \delta)}{2\mu} \right\rfloor$ , where  $\mu$  denotes the Laplacian spectral radius of  $\Gamma$  and  $\delta$  denotes its minimum degree.

*Proof.* Let  $S \subset V$  be an independent set of maximum cardinality  $\alpha(\Gamma)$ . Since the set  $V \setminus S$  is a global offensive alliance in  $\Gamma = (V, E)$ , then

$$\gamma_o(\Gamma) + \alpha(\Gamma) \leq n. \quad (5)$$

If  $|V \setminus S| = 1$ , then  $\Gamma = K_{1, n-1}$  and  $\gamma_o(\Gamma) = 1$ . If  $|V \setminus S| \neq 1$ , let  $V \setminus S = X \cup Y$  be a partition of  $V \setminus S$  such that the edge-cut between  $X$  and  $Y$  has the maximum cardinality. Suppose  $|X| \leq |Y|$ . For every  $v \in Y$ ,  $|N_S(v)| \geq 1$  and  $|N_X(v)| \geq |N_Y(v)|$ . Therefore, the set  $W = S \cup X$  is a global offensive alliance in  $\Gamma$ , i.e., for every  $v \in Y$ ,  $|N_W(v)| \geq |N_Y(v)| + 1$ . Then we have,  $2|X| + \alpha(\Gamma) \leq n$  and  $\gamma_o(\Gamma) \leq |X| + \alpha(\Gamma)$ . Thus,

$$2\gamma_o(\Gamma) - \alpha(\Gamma) \leq n. \quad (6)$$

The bounds i) and ii) follow from (5) and (6).

The proof of *iii*) follows in the spirit of the proof of (6): in this case we take  $S \subset V$  as a dominating set of minimum cardinality. Finally, it was shown in [8] that

$$\alpha(\Gamma) \leq \frac{n(\mu - \delta)}{\mu}.$$

Thus, by (6) we obtain *iv*).  $\square$

The above bounds are attained, for instance, for the cocktail-party graph  $\Gamma = K_6 - F \cong K_{2,2,2}$  where  $n = \mu = 6$ ,  $\delta = 4$ ,  $\alpha(\Gamma) = \gamma(\Gamma) = 2$  and  $\gamma_o(\Gamma) = 4$ .

In the spirit of the proof of *iii*) we obtain

$$2\gamma_o(\Gamma) - \gamma_c \leq n, \quad (7)$$

where  $\gamma_c(\Gamma)$  denotes the connected-domination number of  $\Gamma$ . Moreover, it was shown in [5] that if  $\Gamma$  is a connected graph of order  $n$  and maximum degree  $\Delta$ , then

$$\gamma_c \leq n - \Delta. \quad (8)$$

Thus, by (7) and (8) we obtain

$$\gamma_o(\Gamma) \leq \left\lfloor \frac{2n - \Delta}{2} \right\rfloor. \quad (9)$$

This bound improves *ii*) if  $\Delta > \frac{2n}{3}$ .

**Theorem 2.** *For all connected graph  $\Gamma$  of order  $n$ ,*

$$i) \quad \gamma_{\delta}(\Gamma) \leq \left\lfloor \frac{n + \gamma_2(\Gamma)}{2} \right\rfloor, \text{ where } \gamma_2(\Gamma) \text{ denotes the 2-domination number of } \Gamma.$$

*If the minimum degree of  $\Gamma$  is greatest or equal to two, then*

$$ii) \quad \gamma_{\delta}(\Gamma) \leq n - \alpha(\Gamma), \text{ where } \alpha(\Gamma) \text{ denotes the independence number of } \Gamma;$$

$$iii) \quad \gamma_{\delta}(\Gamma) \leq \left\lfloor \frac{5n}{6} \right\rfloor;$$

$$iv) \quad \text{if } \Gamma \text{ is a cubic graph, then } \gamma_{\delta}(\Gamma) \leq \left\lfloor \frac{3n}{4} \right\rfloor.$$

*Proof.* Let  $H \subset V$  be a 2-dominating set of minimum cardinality. If  $|V \setminus H| = 1$ , then  $\gamma_2(\Gamma) = n - 1$  and  $\gamma_{\delta}(\Gamma) \leq n - 1$ . If  $|V \setminus H| \neq 1$ , let  $V \setminus H = X \cup Y$  be a partition of  $V \setminus H$  such that the edge-cut between  $X$  and  $Y$  has the maximum cardinality. Suppose  $|X| \leq |Y|$ . For every  $v \in Y$ ,  $|N_H(v)| \geq 2$  and  $|N_X(v)| \geq |N_Y(v)|$ . Therefore, the set  $W = H \cup X$  is a global strong offensive alliance in  $\Gamma$ , i.e., for every  $v \in Y$ ,  $|N_W(v)| \geq |N_Y(v)| + 2$ . Then we have,

$$2|X| + \gamma_2(\Gamma) \leq n \quad (10)$$

and

$$\gamma_{\delta}(\Gamma) \leq |X| + \gamma_2(\Gamma). \quad (11)$$

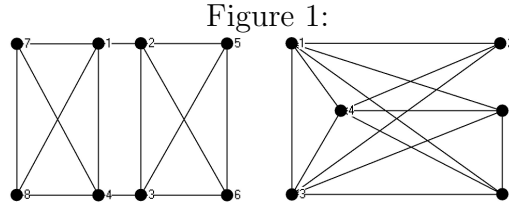
Thus, by (10) and (11), *i*) follows.

Let  $S \subset V$  be an independent set of maximum cardinality  $\alpha(\Gamma)$ . Since  $\delta \geq 2$ , the set  $V \setminus S$  is a global strong offensive alliance in  $\Gamma = (V, E)$ . Hence, *ii*) follows. On the other hand, it was shown in [1] that

$$\delta \geq 2 \Rightarrow \gamma_2(\Gamma) \leq \frac{2n}{3}. \quad (12)$$

So, by *i*) and (12), *iii*) follows.

Finally, if  $\Gamma$  is connected with maximum degree  $\Delta \leq 3$ , then for all global strong offensive alliance  $S$  such that  $|S| = \gamma_{\delta}(\Gamma)$ ,  $V \setminus S$  is an independent set. Thus,  $m \leq 3(n - \gamma_{\delta}(\Gamma)) + \gamma_{\delta}(\Gamma)$ . Hence, the result follows.  $\square$



The bounds *i*) and *ii*) are attained, for instance, for the cocktail-party graph  $\Gamma = K_6 - F$  where  $\gamma_2(\Gamma) = 2$  and  $\gamma_{\delta}(\Gamma) = 4$ . The bound *iii*), is attained, for instance, for the left hand side graph of Figure 1: in this case  $\gamma_{\delta}(\Gamma) = 6$ . Example of equality in *iv*) is  $\Gamma = K_3 \times K_2$ . We emphasize that there are graphs with minimum degree one, such that bounds *ii*) and *iii*) fail. This is, for instance, the case of the star graph,  $\Gamma = K_{1,r}$ , with  $r \geq 6$ . In this case  $n = r + 1$  and  $\gamma_{\delta}(\Gamma) = \alpha(\Gamma) = r$ .

### 3 Bounding below the global offensive alliance number

**Theorem 3.** *For all connected graph  $\Gamma$  of order  $n$ , minimum degree  $\delta$  and maximum degree  $\Delta$ ,*

$$\begin{aligned} i) \quad \gamma_0(\Gamma) &\geq \begin{cases} \left\lceil \frac{n(\delta+1)}{2\Delta+\delta+1} \right\rceil & \text{if } \delta \text{ odd;} \\ \left\lceil \frac{n\delta}{2\Delta+\delta} \right\rceil & \text{otherwise;} \end{cases} \\ ii) \quad \gamma_{\hat{0}}(\Gamma) &\geq \begin{cases} \left\lceil \frac{n(\delta+3)}{2\Delta+\delta+3} \right\rceil & \text{if } \delta \text{ odd;} \\ \left\lceil \frac{n(\delta+2)}{2\Delta+\delta+2} \right\rceil & \text{otherwise.} \end{cases} \end{aligned}$$

*Proof.* Let  $\gamma_k(\Gamma)$  denotes the  $k$ -domination number of  $\Gamma$ . Since all global strong offensive alliance is a  $\left\lceil \frac{\delta+1}{2} \right\rceil$ -dominating set and all global strong offensive alliance is a  $\left\lceil \frac{\delta+2}{2} \right\rceil$ -dominating set,

$$\gamma_{\left\lceil \frac{\delta+1}{2} \right\rceil}(\Gamma) \leq \gamma_0(\Gamma) \quad (13)$$

and

$$\gamma_{\left\lceil \frac{\delta+2}{2} \right\rceil}(\Gamma) \leq \gamma_{\hat{0}}(\Gamma). \quad (14)$$

On the other hand, for all  $k$ -dominating set  $S \subset V$ ,  $k(n - |S|) \leq \Delta|S|$ . Hence,

$$\gamma_k(\Gamma) \geq \left\lceil \frac{kn}{\Delta + k} \right\rceil. \quad (15)$$

Therefore, the result follows.  $\square$

Examples of equality in above theorem are  $\Gamma = K_{3,3}$  and the 3-cube graph.

The following result provides tight bounds on  $\gamma_o(\Gamma)$  and  $\gamma_{\hat{o}}(\Gamma)$  in terms of the order and size of  $\Gamma$ .

**Theorem 4.** *For all graph  $\Gamma$  of order  $n$  and size  $m$ ,*

$$\gamma_o(\Gamma) \geq \left\lceil \frac{3n - \sqrt{9n^2 - 8n - 16m}}{4} \right\rceil$$

and

$$\gamma_{\hat{o}}(\Gamma) \geq \left\lceil \frac{3n + 1 - \sqrt{9n^2 - 10n - 16m + 1}}{4} \right\rceil.$$

*Proof.* If  $S$  denotes a global offensive alliance in  $\Gamma = (V, E)$ , then

$$2m = \sum_{v \in V \setminus S} \delta(v) + \sum_{v \in S} \delta(v) \leq (n - |S|)(2|S| - 1) + |S|(n - 1). \quad (16)$$

Hence, solving  $2|S|^2 - 3n|S| + 2m + n \leq 0$  we obtain the bound on  $\gamma_o(\Gamma)$ . The bound on  $\gamma_{\hat{o}}(\Gamma)$  is derived by analogy by using

$$2m \leq (n - |S|)(2|S| - 2) + |S|(n - 1) \quad (17)$$

instead of (16).  $\square$

Example of equality in the above bounds is the right hand side graph of Figure 1 where  $S = \{2, 6, 5\}$  is a minimal global offensive alliance and  $S' = \{1, 3, 4\}$  is a minimal global strong offensive alliance. Even so, the following bounds, expressed in terms of the order, size, and the maximum degree of  $\Gamma$ , improve the previous result.

**Theorem 5.** *For all graph  $\Gamma$  of order  $n$ , size  $m$  and maximum degree  $\Delta$ ,*

$$\gamma_0(\Gamma) \geq \left\lceil \frac{2m + n}{3\Delta + 1} \right\rceil \quad \text{and} \quad \gamma_{\hat{o}}(\Gamma) \geq \left\lceil \frac{2(m + n)}{3\Delta + 2} \right\rceil.$$

*Proof.* If  $S \subset V$ , then

$$|S|\Delta \geq \sum_{v \in V \setminus S} |N_S(v)|. \quad (18)$$

Moreover, if  $S$  is a global offensive alliance in  $\Gamma$ , then

$$\sum_{v \in V \setminus S} |N_S(v)| \geq \sum_{v \in V \setminus S} |N_{V \setminus S}(v)| + (n - |S|). \quad (19)$$



Thus,

$$\begin{aligned}
2m &= \sum_{v \in V \setminus S} \delta(v) + \sum_{v \in S} \delta(v) \\
&= \sum_{v \in V \setminus S} |N_S(v)| + \sum_{v \in V \setminus S} |N_{V \setminus S}(v)| + \sum_{v \in S} \delta(v) \\
&\leq 2 \sum_{v \in V \setminus S} |N_S(v)| + |S| - n + \sum_{v \in S} \delta(v) \\
&\leq (3\Delta + 1)|S| - n.
\end{aligned}$$

So, the bound on  $\gamma_0(\Gamma)$  follows. If the global offensive alliance  $S$  is strong, then we have

$$\sum_{v \in V \setminus S} |N_S(v)| \geq \sum_{v \in V \setminus S} |N_{V \setminus S}(v)| + 2(n - |S|). \quad (20)$$

Basically, the bound on  $\gamma_\delta(\Gamma)$  follows as before by using (20) instead of (19).  $\square$

The above bounds are reached, for instance, in the case of the 3-cube graph  $\Gamma = K_2 \times K_2 \times K_2$ , where  $\gamma_o(\Gamma) = \gamma_\delta(\Gamma) = 4$ . Notice that Theorem 4 only gives  $\gamma_o(\Gamma) \geq 2$ .

As we can see in [7], we can obtain bounds on the alliance numbers from the spectrum of  $\Gamma$  or from the Laplacian spectrum of  $\Gamma$ . For instance, the following result was proved in [7]. For completeness we include the proof of this result.

**Theorem 6.** *For all graph  $\Gamma$  of order  $n$  and size  $m$ , minimum degree  $\delta$  and Laplacian spectral radius  $\mu$ ,*

$$\gamma_o(\Gamma) \geq \left\lceil \frac{n}{\mu} \left\lceil \frac{\delta + 1}{2} \right\rceil \right\rceil \quad \text{and} \quad \gamma_\delta(\Gamma) \geq \left\lceil \frac{n}{\mu} \left( \left\lceil \frac{\delta}{2} \right\rceil + 1 \right) \right\rceil.$$

*Proof.* It was shown in [3] that the Laplacian spectral radius of  $\Gamma$ ,  $\mu$ , satisfies

$$\mu = 2n \max \left\{ \frac{\sum_{v_i \sim v_j} (w_i - w_j)^2}{\sum_{v_i \in V} \sum_{v_j \in V} (w_i - w_j)^2} : w \neq \alpha \mathbf{j} \text{ for } \alpha \in \mathbb{R} \right\}. \quad (21)$$

Let  $S \subset V$ . From (21), taking  $w \in \mathbb{R}^n$  defined as

$$w_i = \begin{cases} 1 & \text{if } v_i \in S; \\ 0 & \text{otherwise,} \end{cases}$$

we obtain

$$\mu \geq \frac{n \sum_{v \in V \setminus S} |N_S(v)|}{|S|(n - |S|)}. \quad (22)$$

Moreover, if  $S$  is a global offensive alliance in  $\Gamma$ ,

$$|N_S(v)| \geq \left\lceil \frac{\delta(v) + 1}{2} \right\rceil \quad \forall v \in V \setminus S. \quad (23)$$

Thus, (22) and (23) lead to

$$\mu \geq \frac{n}{|S|} \left\lceil \frac{\delta + 1}{2} \right\rceil. \quad (24)$$

Therefore, solving (24) for  $|S|$ , and considering that it is an integer, we obtain the bound on  $\gamma_{ao}(\Gamma)$ . If the global offensive alliance  $S$  is strong, then

$$|N_S(v)| \geq \left\lceil \frac{\delta(v)}{2} \right\rceil + 1 \quad \forall v \in V \setminus S. \quad (25)$$

Thus, (22) and (25) lead to the bound on  $\gamma_{\delta}(\Gamma)$ .  $\square$

If  $\Gamma$  is the Petersen graph, then  $\mu = 5$ . Thus, Theorem 6 leads to  $\gamma_o(\Gamma) \geq 4$  and  $\gamma_{\delta}(\Gamma) \geq 6$ . Therefore, the above bounds are tight.

## 4 Offensive alliances and connected subgraphs

An offensive alliance (global offensive alliance)  $S$  in  $\Gamma$  is *minimal* if no proper subset of  $S$  is an offensive alliance (global offensive alliance) in  $\Gamma$ .

**Theorem 7.** *Let  $\Gamma = (V, E)$  be a connected graph of order  $n$  and diameter  $D(\Gamma)$ . If  $\Gamma$  has a minimal (global) offensive alliance  $S$  such that  $\langle V \setminus S \rangle$  is connected, then  $D(\Gamma) \leq n - |S| + 1$ .*

*Proof.* If  $S \subset V$  is a minimal (global) offensive alliance in  $\Gamma$  then  $V \setminus S$  is a dominating set in  $\Gamma$ . So, if  $\langle V \setminus S \rangle$  is connected, then  $D(\Gamma) \leq D(\langle V \setminus S \rangle) + 2$ . Hence,  $D(\Gamma) \leq n - |S| + 1$ .  $\square$

We remark that there are graphs such that for every minimal (global) offensive alliance  $S$ ,  $\langle V \setminus S \rangle$  is not connected. For instance, the case of the 3-cube graph.

The above bound is tight. Let  $\Gamma$  be the left hand side graph of Figure 2. In this case the set  $S = \{1, 3, 5\}$  is a minimal global offensive alliance and  $V \setminus S = \{2, 4\}$  is connected. Thus,  $3 = D(\Gamma) \leq n - |S| + 1 = 3$ .

**Theorem 8.** *Let  $\Gamma = (V, E)$  be a graph of order  $n$  and maximum degree  $\Delta$ . For all minimal global offensive alliance  $S$  such that  $\langle V \setminus S \rangle$  is connected,*

$$|S| \geq \left\lceil \frac{3n - 2}{\Delta + 3} \right\rceil.$$

*Moreover, for all minimal global strong offensive alliance  $S$  such that  $\langle V \setminus S \rangle$  is connected,*

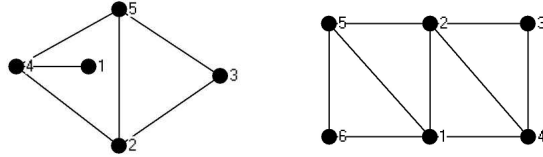
$$|S| \geq \left\lceil \frac{4n - 2}{\Delta + 4} \right\rceil.$$

*Proof.* Let  $S \subset V$ . As  $\langle V \setminus S \rangle$  is connected,

$$\sum_{v \in V \setminus S} |N_{V \setminus S}(v)| \geq 2(n - |S| - 1). \quad (26)$$

So, the first bound follows, by (18), (19) and (26). The second bound is derived by analogy by using (20) instead of (19).  $\square$

Figure 2:



The above bounds are tight. If  $\Gamma$  is the left hand side graph of Figure 2, then  $S = \{1, 3, 5\}$  is a minimal global offensive alliance in  $\Gamma$  and  $V \setminus S = \{2, 4\}$  is connected. Moreover, if  $\Gamma$  is the right hand side graph of Figure 2, then  $S = \{3, 4, 5, 6\}$  is a minimal global strong offensive alliance in  $\Gamma$  and  $V \setminus S = \{1, 2\}$  is connected.

We define the *global-connected offensive alliance number*,  $\gamma_{co}(\Gamma)$ , (respectively, *global-connected strong offensive alliance number*  $\gamma_{c\hat{o}}(\Gamma)$ ) as the minimum cardinality of any global offensive alliance (respectively, global strong offensive alliance) in  $\Gamma$  whose induced subgraph is connected.

**Theorem 9.** *Let  $\Gamma$  be a simple graph of order  $n$ , size  $m$ , diameter  $D$  and maximum degree  $\Delta$ . The global-connected offensive alliance number of  $\Gamma$  is bounded by*

$$\gamma_{co}(\Gamma) \geq \left\lceil \frac{2m + n + 2(D-1)^2}{2n + \Delta + 1} \right\rceil$$

*and the global-connected strong offensive alliance number of  $\Gamma$  is bounded by*

$$\gamma_{c\hat{o}}(\Gamma) \geq \left\lceil \frac{2(m + n + (D-1)^2)}{2n + \Delta + 2} \right\rceil.$$

*Proof.* If  $S$  is a global offensive alliance in  $\Gamma = (V, E)$ , then by (19) we have

$$(|S| - 1)(n - |S|) \geq \sum_{v \in V \setminus S} |N_{V \setminus S}(v)|. \quad (27)$$

Thus,

$$(2|S| - 1)(n - |S|) \geq \sum_{v \in V \setminus S} |N_S(v)| + \sum_{v \in V \setminus S} |N_{V \setminus S}(v)| = \sum_{v \in V \setminus S} \delta(v). \quad (28)$$

Therefore,

$$(2|S| - 1)(n - |S|) + \Delta|S| \geq \sum_{v \in V \setminus S} \delta(v) + \sum_{v \in S} \delta(v) = 2m. \quad (29)$$

On the other hand, if  $S$  is a dominating set and  $\langle S \rangle$  is connected, then  $D(\Gamma) \leq D(\langle S \rangle) + 2$ . So,  $D(\Gamma) \leq |S| + 1$ . Hence,

$$2n|S| - n + |S| + \Delta|S| \geq 2m + 2(D(\Gamma) - 1)^2. \quad (30)$$

Thus, the bound on  $\gamma_{co}(\Gamma)$  follows. Basically the bound on  $\gamma_{c\hat{o}}(\Gamma)$  follows as before by using (20) instead of (19).  $\square$

The above bounds are tight, as we show in the following instance. Let  $\Gamma_{3,t}$  be the graph obtained by joining every vertex of the complete graph  $K_3$  with every vertex of the trivial graph of order  $t \geq 8$ . In such case,  $\gamma_{co}(\Gamma_{3,t}) = \gamma_{c\hat{o}}(\Gamma_{3,t}) = 3$  and Theorem 9 leads to  $\gamma_{co}(\Gamma_{r,t}) \geq 3$  and  $\gamma_{c\hat{o}}(\Gamma_{3,t}) \geq 3$ .

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